

## Series 6 (27.10.2011)

Submission: **15:00, November, 3, 2011**, into the boxes next to the room I-21. Write every solution on the separate paper (format A4)! Don't forget to sign your solutions.

It is not sufficient to answer a single number or yes/not. Answers always have to be justified.

Let us define generalized NFA  $(Q, \Sigma, \delta, Q_I, F)$ , where  $Q, \Sigma, \delta, F$  has the same meaning as for ordinary NFA.  $Q_I \subseteq Q$  is nonempty set of initial states. Computation of an generalized NFA on  $x$  starts with configuration  $(q, x), q \in Q_I$ .

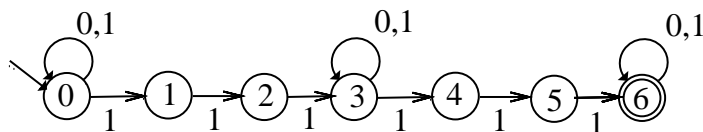
**Exercise 1.** Show that for every generalized NFA  $A'$  there exists equivalent ordinary NFA  $A$ , i.e. such that  $L(A') = L(A)$ .

For the generalized NFA  $A$  we have generalized subset construction of deterministic FA  $\pi(A)$  (which contains only reachable states).

Operation reverse assign to an deterministic FA  $A = (Q, \Sigma, \delta, q_0, F)$  generalized NFA  $\rho(A) = (Q, \Sigma, \rho(\delta), F, \{q_0\})$ , we exchange initial and accepting states and  $q \in \rho(\delta)(q', a) \Leftrightarrow \delta(q, a) = q'$  (arrows are reversed).

**Exercise 2.** Show that  $A' = \pi(\rho(\pi(\rho(A))))$  is deterministic FA which is equivalent with  $A$ , i.e.  $L(A) = L(A')$ . Automaton  $A'$  is minimal, i.e. any other deterministic FA  $B$  such that  $L(B) = L(A)$  has at least so many states as has  $A'$ . (Bonus not compulsory exercise: Prove that  $A'$  is minimal.)

**Exercise 3.** For nondeterministic FA  $M$



create

- using subset construction deterministic FA  $A$ .
- For automaton  $A$  from the part a) construct (and draw)  $\rho(A)$ ,  $\pi(\rho(A))$ ,  $\rho(\pi(\rho(A)))$  a  $\pi(\rho(\pi(\rho(A))))$ .